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Avoiding Traffic Congestion Using Position Based Routing Protocol in VANET

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ABSTRACT— Inter-Vehicle Communication System in a built-up city environment faces greater challenges because of potentially more uneven distribution of vehicular nodes, constrained mobility, and difficult signal reception due to radio obstacles such as high rise buildings. This paper proposes a new position-based routing scheme designed specifically for IVCS in a city environment and for the roadside equipment using a Cognitive Radio to allocate a multiple channel for communicating between vehicles to infrastructure to provide a safety message to avoid a emergency event.

INDEX TERMS— Vehicular Communication; A-star Algorithm; Fuzzy Logic; Cognitive Radio; Roadside Infrastructure; Position Based Routing.

I. INTRODUCTION

A Vehicular Ad-Hoc Network, or VANET, is a form of Mobile Ad-Hoc Network, to provide communication among nearby vehicles and between vehicles and nearest fixed equipment, usually described as Roadside Equipment.

The VANET used to providing safety and comfort for passenger. Having VANET inside vehicle need only small electronic device, which will provide Ad-Hoc Network connectivity for the passengers inside the vehicle. By this device operating this network does not need complicated connection and server communication. Each vehicle equipped with VANET device will be a node in the Ad-Hoc network and can receive and relay others messages through the wireless network

In vehicular Ad-Hoc network using different ad-hoc networking technologies such as Wi-Fi IEEE 802.11 b/g, WiMAX IEEE 802.16, Bluetooth, IRA, ZigBee for easy, accurate, effective and simple communication between vehicles on dynamic mobility.

II. VEHICULAR AD-HOC NETWORKS

VANETs have turned into an important research area over the last few years. VANETs are distinguished from MANET by their hybrid network architectures, node movement characteristics, and new application scenarios.

2.1 CHARACTERISTICS

Drive behavior, constraints on mobility, and high speeds create unique Characteristics in VANETs. These characteristics distinguish them from other mobile ad hoc networks, and the major characteristics are as follows:



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- High mobility and Rapid changing topology
- Geographic position available
- Mobility modeling and predication
- Hard delay constraints
- No power constraint

III. ROUTING INFORMATION USED IN PACKET FORWARDING

This class is divided into two subclasses: topology-based and position-based routing protocols. In topology-based routing, each node should be aware of the network layout, also should able to forward packets using information about available nodes and links in the network. In contrast, position-based routing should be aware of the nodes locations in the packet forwarding.

3.1 TOPOLOGY-BASED ROUTING PROTOCOLS:

Topology-based routing protocol usually a traditional MANET routing protocol, it uses link's information which stored in the routing table as a basis to forward packets from source node to destination; it commonly categorized into three categories (base on underlying architecture) Proactive (periodic), Reactive (on-demand) and Hybrid.

3.1.1 REACTIVE ROUTING PROTOCOLS

Reactive routing protocol (also called as On-Demand) reduce the network overhead by maintaining routes only when needed, that the source node start a process of route discovery and made a route request message to the destination. After the message received by destination it will send a route reply message to the source node.

3.1.1.1 AD-HOC ON-DEMAND DISTANCE VECTOR

AODV offers low network overhead by reducing message flooding in network that when compared to proactive routing protocols, besides reducing the requirement of memory size and by minimizing the routing tables which keep only entries for recent active routes, also keeps next hop for a route rather than the whole route. It also providing dynamically updates for adapting the route condition and eliminated looping in routes by using destination sequence number. However, it causes large delays in a route discovery, also route failure may require a new route discovery which produces additional delays that decrease the data transmission rate and increase the network overhead

3.2 POSITION-BASED ROUTING PROTOCOLS

Position or geographic routing protocol is based on the positional information in routing process; where the source sends a packet to the destination using its geographic position rather than using the network address. When the source need to send a packet, it usually stores the position of the destination in the packet header which will help in forwarding the packet to the destination without needs to route discovery, route maintenance, or even awareness of the network topology.

3.2.1 CHALLENGES OF POSITION-BASED ROUTING IN IVCS

The challenges of position-based routing in a city environment Suppose node s wants to send a packet to node d . Greedy forwarding will fail in this case as there is no neighbor of s , which is nearer to d than s itself. Such a situation is what is commonly known as *local maximum*. Following the strategy in GPSR, the packet enters into perimeter-mode, using the right hand rule to travel through each node on the dotted route, including nodes a , b and c . At b , it is found that c is nearer to d than s , at which the packet enters into perimeter-mode. Thus, the packet switches back to greedy mode at

b , and then reaches its destination d through c . It can be seen that this route is very long in terms of hop count. In fact, s can reach a , and a can reach b , both in one hop.

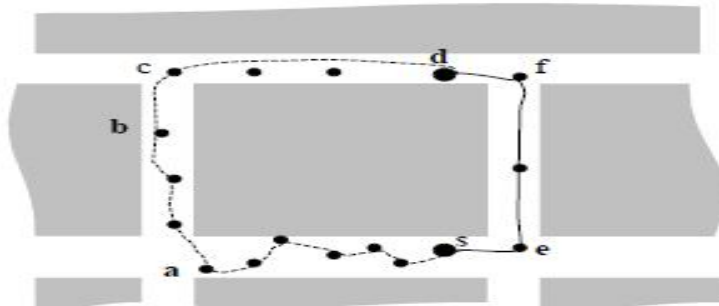


Fig 1 Challenges of position-based routing in IVCS

This shows that the perimeter-mode which packet employs to recover from local maximum is very inefficient and time-consuming. Another observation is that the packet can actually travel from s to d via a route that passes through e and f (shown as solid line), which is much shorter. However, this route is not exploited because the perimeter-mode of GPSR based on right hand rule is biased to a specific direction when selecting for the next hop. It should be noted that in a city environment, the constrained mobility and frequently encountered obstacles can effectively force GPSR to run into perimeter-mode frequently. As a result, the performance of GPSR could deteriorate dramatically, and therefore may not be suitable for IVCS.

3.2.1 ANCHOR-BASED STREET AND TRAFFIC AWARE ROUTING

A-star search algorithm is a widely used graphic searching algorithm. It is also a highly efficient heuristic algorithm used in finding a variable or low cost path. It is considered as one of the best intelligent search algorithms that combines the merits of both depth-first search algorithm and breadth-first algorithm.

A-star path searching algorithm uses the evaluation function (usually denoted $f(n)$) to guide and determine the order in which the search visits nodes in the tree. The evaluation function is given as:

$$f(n) = g(n) + h(n)$$

where $g(n)$ is the actual cost from the initial node (start node) to node n (i.e. the cost finding of optimal path), $h(n)$ is the estimated cost of the optimal path from node n to the target node (destination node), which depends on the heuristic information of the problem area. Generally, A-star algorithm maintains two lists, an OPEN list and a CLOSE list. The OPEN list is a priority queue and keeps track of the nodes in it to find out the next node with least evaluation function to pick. The CLOSE list keeps track of nodes that have already been examined. Initially, the OPEN list contains the starting node. When it iterates once, it takes the top of the priority list, and then checks whether it is the goal node (destination node). If so, the algorithm is done. Otherwise, it calculates the evaluation function of all adjacent nodes and adds them to the OPEN list. After the A-star algorithm is completed, it will find a solution if a solution exists. If it doesn't find a solution, then it can guarantee that no such solution exists. A-star algorithm will find a path with the lowest possible cost. This will depend

heavily upon the quality of the cost function and estimates provided.

A-star algorithm may be expressed as following:

- 1) Put the source node $s_0, f(s_0)$ attached, into the OPEN list. Let the CLOSE list is empty.
- 2) If the OPEN is empty, exit, and the search is fail.
- 3) Move out the first node N form the OPEN list, which has the smallest $f(.)$ in the list, and put it into the CLOSE list; number the node as n .

- 4) If the node N is the goal node, the search is finished, *exits*.
- 5) If the node N cannot spread, turn to *step 2*.
- 6) Spread the node N , there will be a group of nodes, all of which are $f(n)$ attached; add the nodes to the OPEN list, then turn to *step 3*; Especially, for the gotten nodes in this step, some processing will be done as follows:
 - a) Examine the OPEN list and the CLOSE list to find whether (some of) the nodes have been included in them. For the nodes that have been included, if they are ancestor node of the node N , delete them; If they are not (the ancestor node), delete them too, but for they are spread on the second time, it is needed to review them and find whether the corresponding $f(n)$, the back pointers of the nodes and even those of the corresponding descendant nodes are needed to be changed. The rule of such changing is “*choosing the short path based on $f(n)$* .”
 - b) For the nodes that have not been included in the OPEN list and the CLOSE list, put them into the OPEN after assigning the back pointer that points to the node N , then, based on $f(n)$, sort all the nodes in the OPEN list in ascending order.

IV. RELATED WORK

4.1 IMPLEMENTATION OF A-STAR

In the new routing method, the base station prepares the routing schedule and broadcast it to each node. A-star algorithm which is used to find the optimal route from the node to the base station is applied to each node. A-star algorithm creates a tree structure in order to search optimal routing path from a given node to the base station.

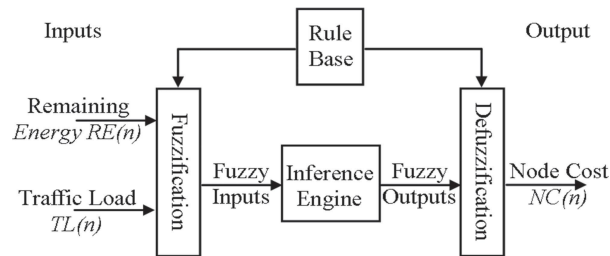


Fig 2. Fuzzy structure with two inputs (remaining energy and traffic load) and one output (node cost).

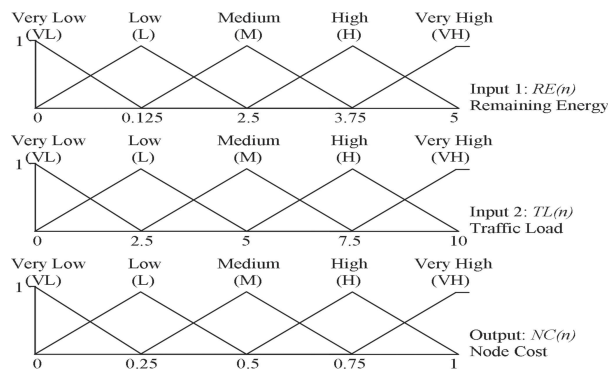


Fig 3 Membership graph for the inputs (remaining energy and traffic load) and the output (node cost).

The tree node is explored based on its *evaluation function $f(n)$* . The function we used is given as:

$$f(n) = NC(n) + (1/MH(n)).$$

Where $NC(n)$ is the node cost of node n , which takes value $[0..1]$, and can be calculated by the fuzzy approach. The fuzzy approach is considered for the remaining energy and the traffic load of node n to calculate the optimal cost for node n . $MH(n)$ is the short distance from node n to the base station. As a result, the node n that has largest $f(n)$ value will be chosen as the optimal node.

4.2 COGNITIVE RADIO

Cognitive radio is used for passing message from roadside unit to vehicles with using a access point with the beacon message for each area to obtain the information of moving vehicles in that area and the vehicle moved out from that region the next access point will get provide the information to the vehicle. So that cognitive radio maintain access point for each and every area to obtain the information of the vehicle and provide the safety message to the vehicle to avoid emergency event occurred in that area so that the driver will slow down the vehicles to avoid accident and traffic in that place. The cognitive radio will have 12channels and it allocate multiple channels to vehicle to provide message to the far vehicle and 8channels will be used at a time for sending message to vehicle and 4channels will be in reserved purpose, if any emergency event occurred remaining 4channels will be allocated for the higher priority vehicles to avoid collision in that region.

V. PERFORMANCE EVALUATION

Energy Model, as implemented in *ns*, is a node attribute. The energy model represents level of energy in a mobile host. The energy model in a node has a initial value which is the level of energy the node has at the beginning of the simulation. This is known as *initialEnergy_*. It also has a given energy usage for every packet it transmits and receives. These are called *txPower_* and *rxPower_*. *EnergyModel(energy)* requires the initial-energy to be passed along as a parameter.

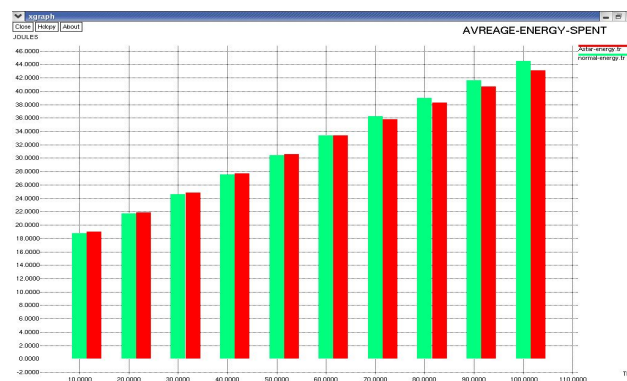


Fig 4 Energy Estimation for AODV and A-STAR

The other class methods are used to decrease the energy level of the node for every packet transmitted (*DecrTxEnergy(txtime, P_tx)*) and every packet received (*DecrRcvEnergy(rcvtime, P_rcv)*) by the node. *P_tx* and *P_rcv* are the transmitting and receiving power (respectively) required by the node's interface or PHY. At the beginning of simulation, *energy_* is set to *initialEnergy_* which is then decremented for every transmission and reception of packets at the node. When the energy level at the node goes down to zero, no more packets can be received or transmitted by the node. If tracing is turned on, line

The performance is made for delay between AODV and A-STAR to provide the delay between the vehicles for the communication using routing protocol with x has a time and y has a milliseconds(ms)

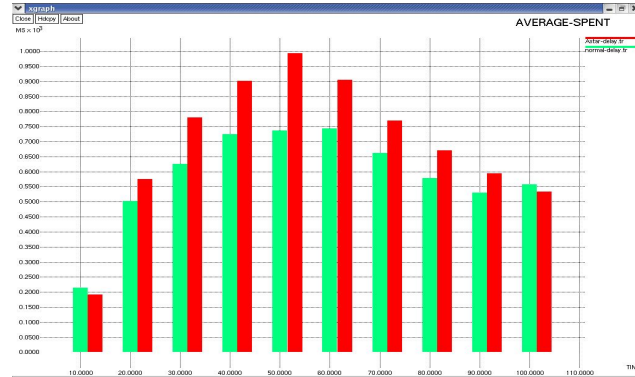


Fig 5 Delay for AODV and A-STAR

The performance evaluation is made to find the throughput for the routing protocol for the communication happening between vehicles and in x-axis contains time and in y-axis contains kilo-bits per seconds(KBPS).



Fig 6 Through-Put for AODV and A-STAR

The performance evaluation has been done for the packet delivery ratio for communication using routing protocol in x-axis time and y-axis contains pdr.

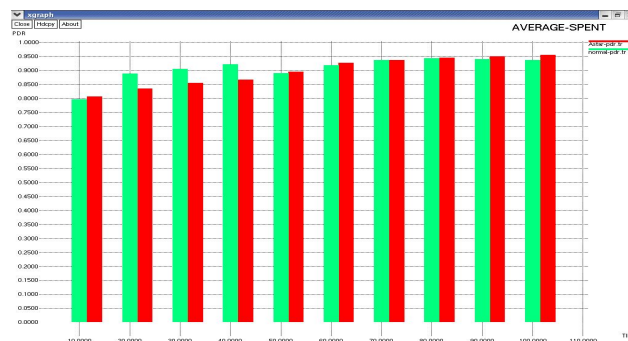


Fig 7 Packet Delivery Ratio for AODV and A-STAR

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This project has high social relevance and usefulness to provide safety message to the far vehicles to avoid traffic congestion in the region using access point with the beacon message to provide the regular update in the network for vehicle to vehicle communication as well as vehicle to infrastructure to pass the message to all the area in the network.

VII. ACKNOWLEDGMENT

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