



# Identifying Malicious Nodes in Wireless Sensor Networks using Node Classification

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**ABSTRACT:** Wireless sensor network is deployed and operated in an unattended and hostile environment to monitor events, produce and transmit data. Nodes in the sensor network could be a gateway, base station, storage node, or querying user. Because of the ease of deployment, the low cost of sensor nodes and the capability of self-organization, sensor network is often used to perform monitoring and data collection tasks. Wireless sensors networks finds its major applications in Military and defence networks. When it is deployed in such an environment, it lacks physical protection and is subject to node compromise. After compromising one or multiple sensor nodes, an adversary may launch various attacks [1] to disrupt the inter-network communication.

**Keywords:** Wireless sensor Networks, Packet Seal, Adversary node, Malicious node

## I. INTRODUCTION

An adversary node could launch an attack of packet dropping to disrupt communication in wireless multi hop sensor networks. Many schemes have been proposed to mitigate and reduce such attacks, but very few can effectively and efficiently identify the intruders for packet drop. A widely used countermeasure is multipath forwarding [2],[3],[4],[5] in which data packets are forwarded in multiple paths and hence packet dropping though not in all paths but could be reduced to a considerable extent but the threat of intruder has not been answered. To address these problems, we propose a simple yet effective scheme, which can identify misbehaving malicious nodes that drop packets by continuously monitoring the behaviours of the nodes in the networks [6], [7], [8], [9], [10], [11].

## II. THE PROPOSED SCHEME

Our proposed scheme contains three techniques

### A. Node Monitoring:

To locate and identify packet droppers, it has been proposed that nodes are continuously monitored for forwarding behaviours and reputation [Bad and suspiciously Bad] of every node is published among the network and maintained in Central node [Sink].

### B. Packet Sealing:

In this scheme, when the sensor data are transmitted by nodes to sink, each packet sender or forwarder seals the data by adding a small number of extra bits called packet seals, from which sink could obtain useful data related to the transmission. Based on the packet seals, the sink can figure out the dropping ratio of every sensor node.



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## C. Node Classification:

The sink identifies and classifies the nodes that are droppers. The behaviour of nodes are traced in variety of scenarios and with the information accumulated in sink, it classifies the nodes as droppers for sure or suspicious droppers.

## III. SYSTEM MODEL

### A. Network assumptions:

The deployment of sensor networks could be such where a large number of sensor nodes are randomly deployed in a two dimensional area. each sensor node generates sensory data periodically and all these nodes collaborate to forward packets containing the data toward a sink. the sink is located within the network. we assume all sensor nodes and the sink are loosely time synchronized [12], which is required by many applications. attack resilient time synchronization schemes, which have been widely investigated in wireless sensor networks [13], [14], can be employed. the sink is aware of the network topology, which can be achieved by requiring nodes to report their neighbouring nodes right after deployment.

Extensive simulation on ns-2 simulator has been conducted to verify the effectiveness and efficiency of the proposed scheme in various scenarios.

### B. Security Assumptions:

The network sink is trustworthy and free of compromise, and the adversary cannot successfully compromise regular sensor nodes during the short and changing topology establishment after the network deployment. this assumption has been widely made in existing work [15], [16].

## IV. IMPLEMENTATION MODEL

In the implementation phase, sensor nodes form a topology which is a directed graph (DG). A routing tree is formed using directed graph. Data flows follow the routing tree structure. In each round, data are transferred through the routing tree to the sink. Each packet sender/forwarder adds a small number of extra bits to the packet (Packet seal) and also encrypts the packet. When one round finishes, based on the extra bits carried in the received packets, the sink runs a node classification algorithm to identify nodes that must be bad (i.e., packet droppers or modifiers) and nodes that are suspiciously bad (i.e., suspected to be packet droppers and modifiers). The routing tree is reshaped every round. As a certain number of rounds have passed, the sink will collect information about node behaviours in different routing topologies. The information includes which nodes are bad for sure, which nodes are suspiciously bad, and the nodes' topological relationship.

The implementation is done in a sequential manner, we first present the algorithm for DG establishment and packet transmission, which is followed by the proposed classification algorithm, tree structure reshaping algorithm.

### A. DG Establishment and Packet Transmission

All sensor nodes form a DG and extract a routing tree from the DG. The sink knows the DG and the routing tree, and shares a unique key with each node. When a node wants to send a packet, it attaches to the packet a sequence number, encrypts the packet only with the key shared with the sink, and then forwards the packet to its parent on the routing tree. When an intermediate node receives a packet, it attaches a few bits to the packet to mark the forwarding path of the packet, encrypts the packet, and then forwards the packet to its parent. On the contrary, a misbehaving intermediate node may drop a packet it receives. On receiving a packet, the sink decrypts it, and thus finds out the original sender and the packet sequence number. The sink tracks the sequence numbers of received packets for every node, and for every certain time interval, which we call a round, it calculates the packet dropping ratio for every node.

Based on the dropping ratio and the knowledge of the topology, the sink identifies packet dropping malicious nodes.

### B. Node Classification Algorithm

In every round, for each sensor node  $u$ , the sink keeps track of the number of packets sent from  $u$ , the sequence numbers of these packets, and the number of flips in the sequence numbers of these packets, In the end of each round, the sink calculates the dropping ratio for each node  $u$ . Suppose  $n_{u, max}$  is the most recently seen sequence number,  $n_{u, flip}$

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is the number of sequence number flips, and  $n_{u,rcv}$  is the number of received packets. The dropping ratio in this round is calculated as follows:

$$d = \frac{nu,flip * Ns + nu,max + 1 - nu,rcv}{nu,flip * Ns + nu,max + 1}$$

Based on the dropping ratio of every sensor node and the tree topology, the sink identifies the nodes that are droppers for sure and that are possibly droppers. After then, for each path from a leaf node to the sink, the nodes' mark pattern in this path can be decomposed into any combination of the following basic patterns, which are also illustrated by Fig. 1:

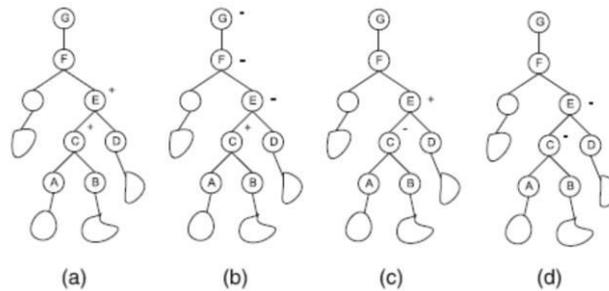


Fig.1. Node Status Pattern

## V. CONCLUSION

The proposed scheme is effective to identify misbehaving forwarders that drop packets. Each packet is encrypted and sealed so as to hide the source of the packet. The packet seal, a small number of extra bits, is added in each packet such that the sink can recover the source of the packet and then figure out the dropping ratio associated with every sensor node. The routing tree structure dynamically changes in each round so behaviours of sensor nodes can be observed in a large variety of scenarios and most of the bad nodes can be identified. Extensive analysis, simulations, and implementation have been conducted and verified the effectiveness of the proposed scheme.

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The proposed work has been simulated in NS-2 and still working to get efficient results in avoiding and detecting intrusion in Wireless sensor networks.

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